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Chempat: Generic Instantiated PQ/T Hybrid Key Encapsulation Mechanisms

Abstract

This document specifies Chempat as a generic family of instantiated Post-Quantum/Traditional (PQ/T) Hybrid Key Exchange Methods (KEMs). The goal is to provide a generic combiner construct that can be analysed separately for security assurance, and to offer concrete instantiated algorithms for integration into protocol and implementations. Identified instances are provided based on traditional Diffie-Hellman key agreement using curves P-256, P-384, X25519, X448, brainpoolP256, brainpoolP384 combined with post quantum methods ML-KEM-768, ML-KEM-1024, Streamlined NTRU Prime sntrup761, and Classic McEliece.

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About This Document

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Status information for this document may be found at <https://datatracker.ietf.org/doc/draft-josefsson-chempat/>.

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Source for this draft and an issue tracker can be found at <https://gitlab.com/jas/ietf-chempat>.

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1. Introduction

To hedge against attacks on a traditional key agreement algorithm such as X25519 [RFC7748] and a post-quantum key encapsulation mechanism (KEM) such as ML-KEM-768 [MLKEM], it is possible to combine both algorithms to derive a shared secret [GHP18] and define the combination mechanism as a new KEM. Using the terminology of [I-D.driscoll-pqt-hybrid-terminology], this combination forms a PQ/T Hybrid Key Encapsulation Mechanism.

Chempat is a generic pattern to create a PQ/T Hybrid Key Encapsulation Mechanism based on at least one post-quantum algorithm and at least one traditional algorithm. The idea is that the Chempat combiner can be analyzed generally and some assurance can be had that it behaves well. For ease of presentation, this document combine one traditional DH-Based KEM algorithm with one post-quantum KEM algorithm.

While a natural approach would be to integrate the generic key combiner construct into protocols and have the protocol and implementation negotiate parameters, that leads to complexity detrimental to security. Therefore this document describe specific instances of Chempat applied on selected algorithms.

2. Motivation

There are many choices that can be made when specifying a hybrid KEM: the constituent KEMs; their security levels; the combiner; and the hash within, to name but a few. Having too many similar options are a burden to the ecosystem.

The above argues for having carefully selected instantiated hybrid KEMs. Each hybrid KEM should be analysed to meet security targets. If that analysis assume specific behaviour of the combiner, or if the analysis become more complex due to the combiner, that leads to more work to re-use the analysis for other combinations. While it would be preferable to only specify one hybrid KEM and analyse that, such as [XWING], cryptographic history suggests that algorithm preferences varies over time.

The argument then is to establish a generic method that can be analysed independent of its component algorithms, such as [KEMCOMBINER]. Generic methods can lead to parametrized protocols and implementations that is more difficult to analyse, and a lack of instantiated algorithm identifiers.

While non-hybrid approaches may eventually be preferable, there are doubts on what properties protocols demand from cryptographic primitives, and some of the properties are different from what have been expected from traditional algorithms [CDM23]. This suggests that some post-quantum KEM's should be used together with a other algorithms to strengthen the properties.

Finally this leads up to our approach to describe a generic method that can be analysed independently of the individual components, with as few parameters as possible in the generic combiner, and to instantiate it with common algorithm choices that make sense for protocols and implementations. That is the essence of Chempat.

3. Comparison to X-Wing

X-Wing [XWING] is a Hybrid PQ/T KEM based on X25519 and ML-KEM-768. Main differences:

- Chempat is applicable to other algorithm combinations, X-Wing's combiner does not extend securely to other KEM combinations.
- Chempat on X25519 with ML-KEM-768 will hash the ML-KEM ciphertext and public key.
- Chempat on X25519 with ML-KEM-768 can provide a per-protocol key-domain separation context string.

4. Comparison to HPKE X25519Kyber768Draft00

HPKE's X25519Kyber768Draft00 [XYBERHPKE] is similar to X-Wing. Main differences to Chempat:

- Chempat is applicable to other algorithm combinations, X25519Kyber768Draft00's combiner does not extend securely to other KEM combinations.
- Chempat hashes the shared secret, to be usable outside of HPKE.
- Chempat hashes the combined ciphertext and public keys.

There is also a different KEM called X25519Kyber768Draft00 [XYBERTLS] which is used in TLS. This one should not be used outside of TLS, as it assumes the presence of the TLS transcript to ensure non malleability.

5. Comparison to KEM Generic Combiner

Chempat is most similar to the generic combiner in [KEMCOMBINER]. Main differences:

- Chempat offers instantiated identified Hybrid KEMs for direct use in protocols and implementations.
- Chempat offers the possibility of a generic simpler security argument for the combiner, whereas [KEMCOMBINER] is parametrized with several algorithm choices and any security analysis needs to be parametrized over the numerous options permitted.
- Chempat has a fixed 32 byte shared secret instead of a variable length shared secret.
- Chempat hashes the public keys of the component KEM's.

6. Design Goals

While Chempat share a lot with [XWING], [XYBERHPKE] and [KEMCOMBINER] the following goals set it apart:

- Allow generic security analysis independent of combinations.
- Provide concrete instantiated algorithm identifiers for several anticipated uses of Hybrid KEM combinations.

We aim for instantiated algorithms of Chempat to be usable for most applications, including specifically HPKE [RFC9180], TLS [RFC8446], OpenPGP [RFC4880] and SSH [RFC4251].

7. Conventions and Definitions

The key words "**MUST**", "**MUST NOT**", "**REQUIRED**", "**SHALL**", "**SHALL NOT**", "**SHOULD**", "**SHOULD NOT**", "**RECOMMENDED**", "**NOT RECOMMENDED**", "**MAY**", and "**OPTIONAL**" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

The following terms are used throughout this document:

string - array of bytes

func1(), func2(a,b) - denote functions called FUNC1 and FUNC2 that takes no parameters and two parameters a and b, respectively.

concat(x0, ..., xN): returns the concatenation of byte strings. concat(0x01, 0x0203, 0x040506) = 0x010203040506.

random(n): return a pseudorandom byte string of length n bytes produced by a cryptographically-secure random number generator.

8. Chempat

Chempat is defined as follows:

```
H = SHA3-256
hybrid_pk = concat(receiver_pk_TKEM, receiver_pk_PQKEM)
hybrid_ct = concat(sender_ct_TKEM, sender_ct_PQKEM)
hybrid_ss = H(concat(ss_TKEM,
                    ss_PQKEM,
                    H(hybrid_ct),
                    H(hybrid_pk),
                    context))
```

The hash function SHA3-256 is defined in [NIST.FIPS.202].

The `hybrid_pk` string is the concatenation of the serialized public-key output from the traditional (`receiver_pk_TEM`) and post-quantum (`receiver_pk_PQKEM`) respectively. To reduce memory usage it is possible to hash the public keys to pre-compute $H(\text{hybrid_pk})$ directly when `hybrid_pk` is received.

The `hybrid_ct` string is the concatenation of the serialized ciphertext output from the traditional (`receiver_ct_TEM`) and post-quantum (`receiver_ct_PQKEM`) respectively. To reduce memory usage it is possible to hash the ciphertext to pre-compute $H(\text{hybrid_ct})$ directly when `hybrid_ct` is received.

The `hybrid_ss` string is the 32-byte output shared secret, formed as the output of the SHA3-256 hash function. The inputs to the hash function is a concatenation of the shared secrets from the traditional (`ss_TKEM`) and post-quantum (`ss_PQKEM`) KEMs with the hashes of the ciphertexts ($H(\text{hybrid_ct})$) and public keys ($H(\text{hybrid_pk})$) together with a variable-length protocol-specific context string.

The context string can be chosen uniquely by the protocol referencing this document. The purpose is to provide protocol domain separation of the generated keys. The content is arbitrary, and in practice the name of the protocol will suffice. Since this results in a new Chempat instance, to reduce combinatorical complexity of parameters, we provide one instance with the context variable set to the name of the Chempat instance, for example "Chempat-X25519-snrup761".

9. Naming

Protocols wishing to utilize a PQ/T Hybrid KEM described in this document **MUST** refer to one of the derived instantiated algorithm identifiers and **MUST NOT** specify a generic construction where the individual algorithms are parameters.

The convention for identifiers is "Chempat-TKEM-PQKEM" replacing "TKEM" and "PQKEM" with a brief mnemonic identifying the traditional and post-quantum algorithm respectively.

10. Use in HPKE

Each Chempat instance satisfy the HPKE KEM interface as follows.

The `SerializePublicKey`, `DeserializePublicKey`, `SerializePrivateKey` and `DeserializePrivateKey` are concatenation and splitting of the known-length component strings.

```

H = SHA3-256

def GenerateKeyPair():
    (pk_T, sk_T) = DHKEM.KeyGen()
    (pk_PQ, sk_PT) = PQKEM.KeyGen()
    return (concat(sk_T, sk_PQ, pk_T, pk_PQ), concat(pk_T, pk_PQ))

# TBA DeriveKeyPair

def Chempat(ss_T, ss_PQ, ct_T, ct_PQ, pk_T, pk_PQ):
    return H(concat(ss_T,
                    ss_PQ,
                    H(concat(ct_T, ct_PQ)),
                    H(concat(pk_T, pk_PQ)),
                    Context))

def Encapsulate(pk):
    pk_T = pk[0:DHKEM.Npk]
    pk_PQ = pk[DHKEM.Npk:PQKEM.Npk-DHKEM.Npk]
    (ss_T, ct_T) = DHKEM.Encap(pk_T)
    (ss_PQ, ct_PQ) = PQKEM.Encap(pk_PQ)
    ss = Chempat(ss_T, ss_PQ, ct_T, ct_PQ, pk_T, pk_PQ)
    ct = concat(ct_T, ct_PQ)
    return (ss, ct)

def Decapsulate(ct, sk):
    ct_T = ct[0:DHKEM.Nenc]
    ct_PQ = ct[DHKEM.Nenc:PQKEM.Nenc-DHKEM.Nenc]
    sk_PQ = sk[0:DHKEM.Nsecret]
    sk_T = sk[DHKEM.Nsecret:PQKEM.Nsecret-DHKEM.Nsecret]
    pk_T = sk[0:DHKEM.Npk]
    pk_PQ = sk[DHKEM.Npk:PQKEM.Npk-DHKEM.Npk]
    ss_T = DHKEM.Decap(ct_T, sk_T)
    ss_PQ = PQKEM.Decap(ct_PQ, sk_PQ)
    return Chempat(ss_T, ss_PQ, ct_T, ct_PQ, pk_T, pk_PQ)

```

Chempat does not provide authenticeted KEMs and does not support AuthEncap() or AuthDecap() of [RFC9180].

Context is a string provided by the protocol referencing this document, or if not provided corresponds to the name of the Chempat instance, such as "Chempat-X25519-sntrup761".

Nsecret is 32 for all Chempat instances, and Nenc, Npk, and Nsk depends on the underlying components.

11. Chempat-X25519-sntrup761

This algorithm is instantiated using the TKEM as DHKEM(X25519, HKDF-SHA256) from [RFC9180] and PQKEM as a HPKE variant of sntrup761 from [NTRUPrimePQCS] [NTRUPrime].

The DHKEM.Nsecret, DHKEM.Nenc, DHKEM.Npk, DHKEM.Nsk are all 32 for X25519 per Section 7.1 of [RFC9180].

The PQKEM.Nsecret is 32, PQKEM.Nenc is 1039, PQKEM.Npk is 1158 and PQKEM.Nsk is 1763 for sntrup761 per [NTRUPrimePQCS].

Thus Nenc is 1071, Npk is 1190 and Nsk is 1795 for Chempat-X25519-sntrup761.

12. Chempat with Classic McEliece with X448 and X25519

This is a set of mechanisms implemented the same way but with different component algorithms and parameter lengths.

This algorithm is instantiated using the TKEM as DHKEM(X, HKDF-SHA512) from [RFC9180] and PQKEM as a HPKE variant of M from [MCELIECE] [CM-spec], substituting X and M for the particular algorithm from the tables below. Sizes for DHKEM for X25519 and X448 as per Section 7.1 of [RFC9180], and sizes for PQKEM as per [CM-spec].

The f and non-f versions are interoperable. The f versions have faster key generation, while the non-f versions have simpler key generation. For example, a key generated with mceliece6688128f can decapsulate ciphertexts that were encapsulated with mceliece6688128, and vice versa. The secret-key sizes (and formats) are the same, the encapsulation functions are the same, and the decapsulation functions are the same. Implementations of this protocol can chose between f and non-f variants, however the name of the hybrid will use the non-f names.

DHKEM variant	Nsecret	Nenc	Npk	Nsk
X25519	32	32	32	32
X448	64	56	56	56

Table 1: X25519/X448 DHKEM size

PQKEM variant	Nsecret	Nenc	Npk	Nsk
mceliece348864	32	96	261120	6492
mceliece460896	32	156	524160	13608
mceliece6688128	32	208	1044992	13932
mceliece6960119	32	194	1047319	13948
mceliece8192128	32	208	1357824	14120

Table 2: Classic McEliece sizes

Names and sizes of the Chempat hybrids are per table below.

Variant	Nenc	Npk	Nsk
Chempat-X25519-mceliece348864	128	261152	6524
Chempat-X25519-mceliece460896	188	524192	13640

Variant	Nenc	Npk	Nsk
Chempat-X25519-mceliece6688128	240	1045024	13964
Chempat-X25519-mceliece6960119	226	1047351	13980
Chempat-X25519-mceliece8192128	240	1357856	14152
Chempat-X448-mceliece348864	160	261176	6548
Chempat-X448-mceliece460896	220	524216	13664
Chempat-X448-mceliece6688128	272	1045048	13988
Chempat-X448-mceliece6960119	258	1047375	14004
Chempat-X448-mceliece8192128	272	1357880	14176

Table 3: Classic McEliece with X25519/X448

13. Chempat-X25519-ML-KEM-768

This algorithm is instantiated using the TKEM as DHKEM(X25519, HKDF-SHA256) from [RFC9180] and PQKEM as a HPKE variant of ML-KEM-768 from [MLKEM].

Protocols and implementation **MAY** consider [XWING] instead of Chempat-X25519-ML-KEM-768, and the definition of Chempat-X25519-ML-KEM-768 is here for situations when some property of X-Wing is not wanted. Informally and non-conclusively, X-Wing offers better performance and Chempat-X25519-ML-KEM-768 offers re-use of the generic security claims on Chempat and a per-protocol key-separation context string.

The DHKEM.Nsecret, DHKEM.Nenc, DHKEM.Npk, DHKEM.Nsk are all 32 for X25519 per Section 7.1 of [RFC9180].

The PQKEM.Nsecret is 32, PQKEM.Nenc is 1088, PQKEM.Npk is 1184 and PQKEM.Nsk is 2400 for ML-KEM-768 per [MLKEM].

Thus Nenc is 1120, Npk is 1216 and Nsk is 2432 for Chempat-X25519-ML-KEM-768.

14. Chempat-X448-ML-KEM-1024

This algorithm is instantiated using the TKEM as DHKEM(X448, HKDF-SHA512) from [RFC9180] and PQKEM as a HPKE variant of ML-KEM-1024 from [MLKEM].

For X448 DHKEM.Nsecret is 64, DHKEM.Nenc is 56, DHKEM.Npk is 56, DHKEM.Nsk is 56 per Section 7.1 of [RFC9180].

The PQKEM.Nsecret is 32, PQKEM.Nenc is 864, PQKEM.Npk is 1568 and PQKEM.Nsk is 2400 for ML-KEM-1024 per [MLKEM].

Thus Nenc is 1120, Npk is 1624 and Nsk is 2456 for Chempat-X25519-ML-KEM-1024.

15. Chempat-P256-ML-KEM-768

This algorithm is instantiated using the TKEM as DHKEM(P-256, HKDF-SHA256) from [RFC9180] and PQKEM as a HPKE variant of ML-KEM-768 from [MLKEM].

For P256 DHKEM.Nsecret is 32, DHKEM.Nenc is 65, DHKEM.Npk is 65, DHKEM.Nsk is 32 per Section 7.1 of [RFC9180].

The PQKEM.Nsecret is 32, PQKEM.Nenc is 1088, PQKEM.Npk is 1184 and PQKEM.Nsk is 2400 for ML-KEM-768 per [MLKEM].

Thus Nenc is 1153, Npk is 1249 and Nsk is 2432 for Chempat-P256-ML-KEM-768.

16. Chempat-P384-ML-KEM-1024

This algorithm is instantiated using the TKEM as DHKEM(P-384, HKDF-SHA384) from [RFC9180] and PQKEM as a HPKE variant of ML-KEM-1024 from [MLKEM].

For P384 DHKEM.Nsecret is 48, DHKEM.Nenc is 97, DHKEM.Npk is 97, DHKEM.Nsk is 48 per Section 7.1 of [RFC9180].

The PQKEM.Nsecret is 32, PQKEM.Nenc is 864, PQKEM.Npk is 1568 and PQKEM.Nsk is 2400 for ML-KEM-1024 per [MLKEM].

Thus Nenc is 961, Npk is 1665 and Nsk is 2448 for Chempat-P384-ML-KEM-1024.

17. Chempat-brainpoolP256-ML-KEM-768

This algorithm is instantiated using the TKEM as DHKEM(brainpoolP256, HKDF-SHA256) from [RFC9180] [RFC5639] and PQKEM as a HPKE variant of ML-KEM-768 from [MLKEM].

For brainpoolP256 DHKEM.Nsecret is 32, DHKEM.Nenc is 65, DHKEM.Npk is 65, DHKEM.Nsk is 32.

The PQKEM.Nsecret is 32, PQKEM.Nenc is 1088, PQKEM.Npk is 1184 and PQKEM.Nsk is 2400 for ML-KEM-768 per [MLKEM].

Thus Nenc is 1153, Npk is 1249 and Nsk is 2432 for Chempat-brainpoolP256-ML-KEM-768.

18. Chempat-brainpoolP384-ML-KEM-1024

This algorithm is instantiated using the TKEM as DHKEM(brainpoolP384, HKDF-SHA384) from [RFC9180] [RFC5639] and PQKEM as a HPKE variant of ML-KEM-1024 from [MLKEM].

For brainpoolP384 DHKEM.Nsecret is 48, DHKEM.Nenc is 97, DHKEM.Npk is 97, DHKEM.Nsk is 48. The PQKEM.Nsecret is 32, PQKEM.Nenc is 864, PQKEM.Npk is 1568 and PQKEM.Nsk is 2400 for ML-KEM-1024 per [MLKEM].

Thus Nenc is 961, Npk is 1665 and Nsk is 2448 for Chempat-brainpoolP384-ML-KEM-1024.

19. Security Considerations

Chempat is intended to be secure if SHA3 is secure and either the traditional algorithm is secure or the post-quantum algorithm is secure.

The security considerations of each component algorithm are inherited.

Cryptographic algorithms and parameters will be broken or weakened over time. Blindly implementing supported groups listed here is not advised. Implementers and users need to check that the cryptographic algorithms listed continue to provide the expected level of security.

20. IANA Considerations

Protocols that provide a Context variable will need to register their own key-domain separate identifiers. The registrations below are when Chempat instances are used with their default value of Context.

This document requests/registers new entries to the "HPKE KEM Identifiers" registry as follows.

Value	KEM	Nsecret	Nenc	Npk	Nsk	Auth	Reference
TBD	Chempat-X25519-sntrup761	32	1071	1190	1795	No	THISRFC
TBD	Chempat-X25519-mceliece348864	32	128	261152	6524	No	THISRFC
TBD	Chempat-X25519-mceliece460896	32	188	524192	13640	No	THISRFC
TBD	Chempat-X25519-mceliece6688128	32	240	1045024	13964	No	THISRFC
TBD	Chempat-X25519-mceliece6960119	32	226	1047351	13980	No	THISRFC
TBD	Chempat-X25519-mceliece8192128	32	240	1357856	14152	No	THISRFC
TBD	Chempat-X448-mceliece348864	32	160	261176	6548	No	THISRFC

Value	KEM	Nsecret	Nenc	Npk	Nsk	Auth	Reference
TBD	Chempat-X448-mceliece460896	32	220	524216	13664	No	THISRFC
TBD	Chempat-X448-mceliece6688128	32	272	1045048	13988	No	THISRFC
TBD	Chempat-X448-mceliece6960119	32	258	1047375	14004	No	THISRFC
TBD	Chempat-X448-mceliece8192128	32	272	1357880	14176	No	THISRFC
TBD	Chempat-X25519-ML-KEM-768	32	1120	1216	2432	No	THISRFC
TBD	Chempat-X448-ML-KEM-1024	32	1120	1624	2456	No	THISRFC
TBD	Chempat-P256-ML-KEM-768	32	1153	1249	2432	No	THISRFC
TBD	Chempat-P384-ML-KEM-1024	32	961	1665	2448	No	THISRFC
TBD	Chempat-brainpoolP256-ML-KEM-768	32	1153	1249	2432	No	THISRFC
TBD	Chempat-brainpoolP384-ML-KEM-1024	32	961	1665	2448	No	THISRFC

Table 4: Chempat HPKE KEM Identifiers

This document requests/registers a new entry to the TLS Supported Group registry as follows.

Value	Description	DTLS-OK	Recommended	Reference	Comment
TBD	Chempat-X25519-sntrup761	Y	Y	THISRFC	PQ/T hybrid of X25519 and sntrup761

Table 5: Chempat TLS Supported Groups

21. Acknowledgments

The combiner function was suggested by Daniel J. Bernstein. The document re-use ideas and some text from [XWING], [KEMCOMBINER], [XYBERHPKE] and [RFC9180].

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